

Maximum augmenting sequence of spanning trees

Let G be an unempty simple undirected connected weighted graph, $G = (V, E, w)$, $w: E \rightarrow \{0, 1, 2, 3, \dots\}$. In this problem, we suppose that the weights of all edges of G are mutually different.

At first, we describe the problem less formally. Let T_1 be a spanning tree of G . The central part of the problem is to create another spanning tree of G , say T_2 , which differs a little from T_1 and which we will call an **augment of T_1** . The process of creating T_2 consists of the following:

1. Find a cheap edge e_1 in T_1 .
2. Find an expensive edge e_2 , which belongs to G but not to T_1 .
3. Remove the cheap edge e_1 from T_1 and insert the expensive edge e_2 into T_1 .

Steps 1.– 3. must be completed in such way that the resulting subgraph, now called T_2 , will again be a spanning tree of G . Also, the weight of T_2 must be bigger than the weight of T_1 and maximum possible. As there can be more pairs of edges (e_1, e_2) which satisfy these conditions, we additionally demand that the weight $w(e_1)$ must also be maximum possible among those pairs of edges.

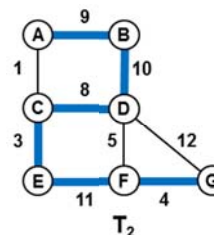
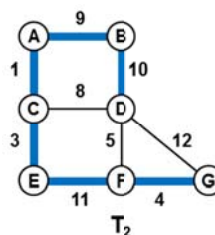
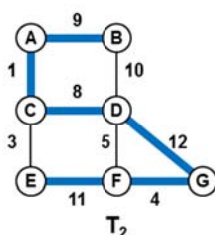
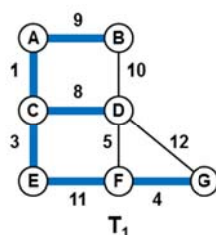
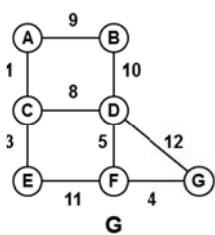


Image 1a. Graph G and its spanning trees T_1 and T_2 . T_2 is an augment of T_1 .

Image 1b. T_2 is not an augment of T_1 .

Image 1c. T_2 is not an augment of T_1 .

Example. In the image above there is a highlighted spanning tree T_1 of the graph G . To create a spanning tree T_2 (also highlighted) which is an augment of T_1 we chose the pair of edges $(\{C, E\}, \{D, G\})$. The weight of T_2 is then 45. This can be seen in the image 1a. We cannot choose, for example, the pair $(\{C, D\}, \{B, D\})$ because the weight of T_2 would be only 38 as can be seen in the image 1b. Also, as illustrated in the image 1c, we cannot choose the pair $(\{A, C\}, \{B, D\})$ because although the weight of T_2 would be 45, which is maximum possible, the weight of the edge $\{A, C\}$ is less than the weight of the first edge of the pair $(\{C, E\}, \{D, G\})$ used in the image 1a.

For the sake of completeness we provide also a formal definition of the augment tree. You may skip this section if the informal description is sufficient for you.

Let T be a spanning tree of G . We say that an ordered pair of edges $(e_1, e_2) \in E \times E$ is a **T -augmenting pair of edges** if both following conditions hold:

1. $e_1 \in E(T)$, $e_2 \in E(G) - E(T)$.
2. The difference $w(e_2) - w(e_1)$ is positive and maximum possible.

We say that a T -augmenting pair of edges (e_1, e_2) is a **proper T -augmenting pair** if the value of $w(e_1)$ is maximum among all T -augmenting pairs of edges.

Let T_1 and T_2 be two spanning trees of a G . We say that T_2 is an **augment of T_1** if there is a proper T_1 -augmenting pair (e_1, e_2) such that $E(T_2) = \{E(T_1) - e_1\} \cup e_2$.

We say that a sequence (T_1, T_2, \dots, T_D) , $(D > 0)$ of spanning trees of G is a **maximum augmenting sequence of G** when both following conditions hold:

1. T_1 is a minimum spanning tree of G .
2. T_k is an augment of T_{k-1} , for $k = 2, 3, \dots, D$.

Note that for some values of D the maximum augmenting sequence might be undefined.

The task

We have to find the weight of the last element of the maximum augmenting sequence of the given graph G . We consider the weight of a graph to be the sum of weights of all its edges.

Input

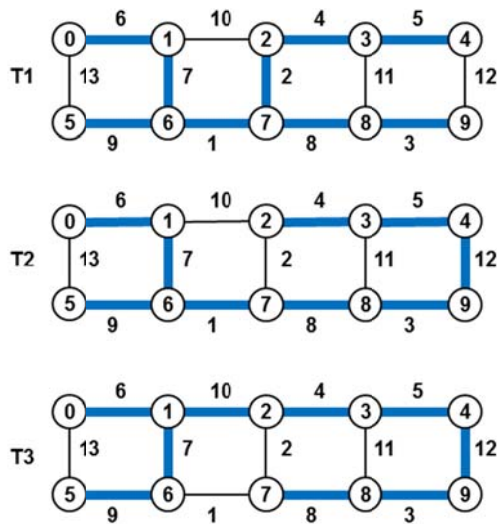
The first line of input contains three positive integers N, M, D separated by space. The integers represent the number of vertices, the number of edges and the value D specified above. Next, there are M lines, each specifies one edge by three integers a, b, c separated by space. The integers a and b represent the edge $\{a, b\}$ and c represent its weight. We suppose that the vertices of the graph are labeled $0, 1, \dots, N-1$. Input values satisfy $D \leq N \leq 2000$.

Output

The output consists of a single line containing the weight of the last element of the maximum augmenting sequence (T_1, T_2, \dots, T_D) of the input graph. The sequence is always defined.

Example 1

Input
 10 13 3
 0 1 6
 1 2 10
 2 3 4
 3 4 5
 5 6 9
 6 7 1
 7 8 8
 8 9 3
 0 5 13
 1 6 7
 2 7 2
 3 8 11
 4 9 12



Output
 64

Image 2. The picture shows the maximum augmenting sequence (T_1, T_2, T_3) of the input graph in Example 1. The edges of the spanning trees are highlighted.

Example 2

Input
 7 21 7
 2 0 0
 3 2 3
 4 2 5
 5 0 7
 6 4 9
 4 1 11
 3 0 13
 4 3 15
 3 1 17
 5 1 19
 1 0 21
 2 1 23
 4 0 25
 6 5 27
 5 4 29
 6 3 31
 5 2 33
 5 3 35
 6 2 37
 6 1 39
 6 0 41

Output
 214