COLT tries to explain why and when machine learning works.

It studies two aspects of machine learning to provide insights for the design of learning algorithms.

- Statistical: how much data is needed to learn good models?
- Algorithmic: how computationally hard is it to learn such models?

COLT usually assumes a simple learning scenario called *concept learning*, which is (roughly) noise-free binary classification learning.

More complex scenarios often have concept learning at their heart.

#### **Concept Learning Elements**

- Instance space: a set X. Elements  $x \in X$  are instances.
- *Concept*: a subset  $C \subseteq X$ .

The algorithm should learn to decide whether  $x \in C$  for any given  $x \in X$ .

Example: X = animals described as tuples of binary variables

	aquatic	airborne	backbone
<i>x</i> =	0	1	0

C =all mammals.

• Learning examples: the learner must get some instances  $x \in X$  with the information whether  $x \in C$  or not.

In general, there is an un-countable number of concepts on X if X is infinite, e.g. X = N.

In practice, we often have prior knowledge about the concepts we want to learn: they belong to a restricted set of concepts

$$\mathcal{C} \subset 2^X$$

which is called a *concept class*.

COLT studies the behavior of learners with respect to selected concept classes.

# Hypothesis Class

A *finite* description of a learner's decision model is called a *hypothesis*. It is an *algorithm* implementing the function  $c : X \to \{0, 1\}$ 

$$c(x) = \begin{cases} 1 \text{ if } x \in C \\ 0 \text{ if } x \notin C \end{cases}$$

Note: in general, there are fewer hypotheses than concepts due to the finiteness of the former.

Learners use constrained languages (rules, polynomials, graphs,  $\dots$ ) to encode their hypotheses. So the algorithmic semantics is implicit. For example, the hypothesis

 $man \land married$ 

which is a *logical conjunction* defines the 'bachelor' concept.

The set of all hypotheses a learner can express is called its hypothesis class.

#### A Continuous-Domain Example

• Instance space  $X = R^2$ 

• Possible concept class C: disks  $(x_1 - a)^2 + (x_2 - b)^2 < r$ 



## A Continuous-Domain Example (cont'd)

- Possible hypothesis class  $\mathcal{H}$ : half-planes  $x_2 ax_1 > b$
- Hypothesis description: (*a*, *b*) (with finite precision number repr.)



## A Continuous-Domain Example (cont'd)

- Possible hypothesis class  $\mathcal{H}$ : neural networks
- Hypothesis description: graph + weights



Instances and hypotheses in continuous domains are largely the topic of a parallel course (Statistical Machine Learning).

Here we focus mainly on discrete domains that allow convenient symbolic representations. Typically:

- Instance attributes are Boolean values;
- Hypotheses are logical formulas (or structures that can be converted to them).

Symbolic representations have the advantage of *understandability* to a human. Important e.g. in medical applications.

Currently studied in the field of "Explainable Al".

A *learning model* is an abstract description of real-life machine-learning scenarios. It defines

- The learner-environment interaction protocol
- How learning examples are conveyed to the learner
- What properties the examples must posses
- What it means to learn successfully

We will discuss two learning models:

- Mistake Bound Learning
- Probably Approximately Correct Learning.

Sometimes, hypotheses are also called models but here we mean a model of learning.

A very simple model assuming an *online* interaction: a concept C is chosen from a fixed concept class and the following is then repeated indefinitely:

- The learner receives an example  $x \in X$
- **2** It predicts whether x is positive  $(x \in C)$  or negative  $(x \notin C)$
- It is told the correct answer (so it can adapt after a wrong prediction)

To define the model, we assume there is a measure n of *instance complexity*. When X consists of fixed-arity tuples, we set n = their arity.

Denote poly(n) to mean "at most polynomial in n". In math expressions,  $f(n) \le poly(n)$  means that f(n) grows at most polynomially. We say that an algorithm *learns concept class* C if for any  $C \in C$ , the number of mistakes it makes is poly(*n*); if such an algorithm exists, C is called *learnable* in the mistake bound model. We will omit "in the mistake bound model" in this section.

Note that the learner

- cannot assume anything about the choice of examples (no i.i.d. or order assumption etc.);
- $\bullet\,$  which learns  ${\mathcal C}\,$  stops making mistakes after a finite number of decisions.

If an algorithm learns C and the maximum time it uses to process a single example is also poly(*n*), we say it learns C *efficiently* and we call C *efficiently* learnable.

Assume  $X = \{0, 1\}^n$   $(n \in N)$  and C consists of all concepts expressible via conjunctions on n variables. Consider the following *generalization* algorithm.

- Initial hypothesis  $h = h_1 \overline{h_1} h_2 \overline{h_2} \dots h_n \overline{h_n}$
- **2** Receive example x, decide "yes" iff h true for  $x (x \models h)$
- If decision was "no" and was wrong, remove all h's literals false for x
- If decision was "yes" and was wrong, output "Concept cannot be described by a conjunction."
- 5 Go to 2

To adapt this algo for  $C = monotone \ conjunctions$  (conj. with no negations), use  $h = h_1 h_2 \dots h_n$  in Step 1.

Let  $C \in C$  be the concept used to generate the examples and c the conjunction that encodes it. Observe and explain why:

- Initial h tautologically false, n literals get deleted from it on first mistake on a positive (in-concept) example, resulting in |h| = n.
- If a literal is in c, it is never deleted from h, so  $c \subseteq h$  (literal-wise).
- At least one literal is deleted on each mistake.
- So the max number of mistakes is  $n + 1 \le poly(n)$ .

Thus the algorithm learns conjunctions (in the MB model) and does so efficiently (time per example is linear in n).

#### So conjunctions are efficiently learnable.

Efficient learnability of conjunctions implies the same for *disjunctions*.

If disjunction *c* defines concept *C* then  $\overline{c}$  is a *conjunction* defining the *complementary* concept  $X \setminus C$ .

Use any efficient conjunction learner to learn  $X \setminus C$ , so the correct answers provided to the learner are according to  $\overline{c}$ .

Then negate the hypothesis returned by the algorithm, obtaining a disjunction for C.

k-CNF (DNF) is the class of CNF (DNF) formulas whose clauses (terms) have at most k literals. For example, 3-CNF includes

 $(a \lor b)(b \lor \overline{c} \lor d)$ 

*k*-CNF is efficiently learnable.

With *n* variables, there are  $n' = \sum_{i=1}^{k} {n \choose i} 2^i \le \text{poly}(n)$  different clauses.

Introduce a new variable for each of the n' clauses and use an efficient learner to learn a monotone conjunction on these variables. Then plug the original clauses for the variables in the resulting conjunction, obtaining a k-CNF formula. This is efficient due to  $n' \leq poly(n)$ .

Analogically, also *k-DNF is efficiently learnable*.

*k*-term DNF (*k*-clause CNF): at most *k* terms (clauses).

No algorithm known for efficient learning of k-term DNF using k-term DNF as the hypothesis class. Same for k-clause CNF.

But k-term  $DNF \subseteq k$ -CNF since any k-term DNF can be written as an equivalent k-CNF by "multiplying-out." E.g.,

 $(abc) \lor (de) \models (a \lor d)(a \lor e)(b \lor d)(b \lor e)(c \lor d)(c \lor e)$ 

So k-term DNF is efficiently learnable by an algorithm using k-CNF as its hypothesis class. This is called *improper* learning.

Analogically: k-clause CNF learnable using k-DNF.